

REMARKS

By this amendment, no claims have been cancelled, Claim 37 has been amended to correct a typographical error that was inadvertently introduced, and new Claims 45-49 have been added. Hence, Claims 8-30, 32-35, 37-40, and 42-49 are pending in the application.

Applicants acknowledge that Claims 16-30, 33-35, 38-40, and 43-44 are in condition for allowance as indicated in the Advisory Action mailed December 1, 2006.

NEW CLAIMS 45-49 ARE ALLOWABLE OVER THE CITED ART

New Claims 45-49 are apparatus claims that feature elements similar to those recited by method Claims 20-24, which are currently in condition for allowance. As a result, for at least the reasons in which Claims 20-24 are allowable over the cited art, Applicants respectfully submit that new Claims 45-49 are also allowable over the cited art and are in condition for allowance.

CLAIMS 8-30, 32-35, 37-40, AND 42-44 ARE ALLOWABLE OVER THE CITED ART

Claims 16-30, 33-35, 38-40, and 43-44 are in condition for allowance as indicated in the Advisory Action mailed December 1, 2006.

Claims 8-15, 32, 37, and 42 stand rejected under 35 U.S.C. § 103(a) as being anticipated by U.S. Patent No. 6,253,236 issued to Troxel et al. ("*Troxel*") in view of U.S. Patent No. 6,493,804 issued to Soltis et al. ("*Soltis*") in further view of U.S. Patent No. 5,913,213 issued to Wikstrom ("*Wikstrom*"). The rejections are traversed.

CLAIM 8

Claim 8 features:

A method of controlling use by concurrent users of a distributed resource on a network, wherein use of the resource is limited to a specified maximum number of concurrent users, the method comprising the computer-implemented steps of: **providing a distributed lock manager process comprising a plurality of local lock manager processes executing on a corresponding plurality of hosts,** **wherein each of the plurality of local lock manager processes may grant a lock on the same resource;** **associating a user identification for each user with one host of the plurality of hosts; and** **responding to a request for the resource associated with a first user having a first user identification associated with a first host of the plurality of hosts by requesting a lock from a first local lock manager process executing on the first host. (emphasis added).**

Even if *Troxel*, *Soltis*, and *Wikstrom* were to be properly combined, at least the above-bolded elements of Claim 8 would still not be disclosed, taught, or suggested by *Troxel* or *Soltis* or *Wikstrom*, either individually or in combination.

Claim 8 recites controlling use by concurrent users of a distributed resource on a network. The resource is limited to a specified maximum number of concurrent users. A distributed lock manager process, comprising a plurality of local lock manager processes, executes on a corresponding plurality of hosts. Each of the plurality of local lock manager processes may grant a lock on the same resource. A user identification for each user is associated with one host of the plurality of hosts. A request, associated with a first user having a first user identification associated with a first host of the plurality of hosts, for the resource is responded to by requesting a lock from a first local lock manager process executing on the first host. By responding to requests for resources in this fashion, the number of concurrent users of a distributed resource on a network may be controlled in a manner that scales to support a large number of clients and avoids a single point of failure.

Troxel, on the other hand, describes enabling a client computer system to access data on a mainframe host computer system in a manner that prevents conflicts for data stored under different types of file access methods and provides security at a file level and recovery from interruption of communication between the host and client computer systems (Col. 2, line 40 – Col. 3, line 21). *Troxel* maintains a concurrent user counter that is incremented each time a client computer system logs into the host computer system to enforce limits on the number of concurrent users that can access the host computer system (See concurrent user counter 354 at Col. 5, lines 26-30).

Troxel lacks any suggestion of a **distributed lock manager process** comprising a **plurality of local lock manager processes**. Instead, the cited portion of *Troxel* teaches a single, centralized component, on the mainframe host computer system, that is responsible for locking resources upon receipt of a client request (Col. 1, lines 15-36). Further, *Troxel* lacks any suggestion of associating a user identification for each user with one host of a plurality of hosts. Also, the cited portion of *Troxel* lacks any suggestion of responding to requests for resources as claimed. In recognition of the deficiencies of *Troxel*'s teachings, the Office Action acknowledges (see page 3):

Troxel does not explicitly teach providing a distributed lock manager process comprising a plurality of local lock manager processes executing on a corresponding plurality of hosts; associating a user identification for each user with one host of the plurality of hosts; and responding to a request for the resource associated with a first user having a first user identification associated with a first host of the plurality of hosts by requesting a lock from a first local lock manager process on the first host.

Thus, the Office Action acknowledges that *Troxel* does not teach or suggest ANY of the above-bolded elements of Claim 8.

The Office Action also cites *Soltis* in rejecting Claim 8. *Soltis* teaches an approach for a global file system that comprises a plurality of clients which each may access one or more data storage devices of the global file system (See FIG. 1; Col. 5, lines 25-45; Col. 10, lines 9-11). The data storage devices may be pooled together into a shared disk memory in a Network Storage Pool (NSP) arrangement (Col. 10, lines 19-23). A data storage device includes one or more locks that are each associated with the use of a storage block of the data storage device. Clients access a particular storage block, of a particular data storage device, by requesting a lock associated with the particular storage block from the particular storage device (see Abstract; Col. 2, line 45 - Col. 3, line 64).

In the approach of *Soltis*, the Network Storage Pool (NSP) provides shared storage resources that are substantially equally accessible to each client in the system (Col. 10, lines 33-35). *Soltis* teaches that each data storage device maintains a set of locks that may be granted to any of the clients accessing data blocks of only that data storage device (see FIG. 7; Col. 13, line 65 – Col. Col. line 32). **A particular data storage device of *Soltis* may only grant locks on the data block of that particular data storage device.** Thus, *Soltis* expressly teaches away from a distributed lock manager process comprising a plurality of local lock manger processes, executing on a corresponding plurality of hosts, that each may grant a lock on the same resource.

As a result, *Soltis* cannot teach a plurality of local lock manager processes that may each grant a lock on the same resource.

Perhaps in acknowledgement of this deficiency in *Soltis*, the Office Action cites a third reference, namely *Wikstrom*, to show “wherein each of the plurality of local lock manager

processes may grant a lock on the same resource.” However, *Wikstrom* also fails to teach or suggest this feature. Instead, *Wikstrom* describes only a single lock manager residing at a node, and each lock manager may only grant locks on resources that reside on the same node as the lock manager.

Wikstrom describes locking **replicated** data objects (see title). The fact that *Wikstrom* describes locking replicated data objects **teaches away** from the approach of Claim 8. In *Wikstrom*, each node of a plurality of nodes has its own lock manager. For example, FIG. 1 of *Wikstrom* shows node 30A comprises lock manager 73A and node 30B comprises lock manger 73B. In *Wikstrom*, each node has its own version of a replicated data object. For example, *Wikstrom* states:

Each node 30 has its own version of a data object X. Specifically, node 30A has hard disk 32A whereon its version of data object X, referenced as X-A, is stored. Similarly, node 30B has hard disk 32B whereon its version of data object X, referenced as X-B, is stored (see Col. 3, lines 38-43).

Further, in *Wikstrom*, each lock manager operates on the version of a replicated data object that resides at the same node as the lock manager. For example, *Wikstrom* states:

Since versions of data object X are stored both at nodes 30A and 30B, when one node updates the value of data object X, the updated value is communicated to the other node so that the other node can likewise have the updated value, thereby maintaining a coordination of the value of the data object X. (see Col. 3, lines 47-53).

As illustrated in FIG. 2, object lock table 220 has a record 220 for each data object pertinent to node A. (see Col. 4, lines 42-43. This teaching shows that the lock table for node A only maintains records for resources at node A).

Consequently, in *Wikstrom*, a single lock manager resides at a node, and each lock manager may only grant locks on resources that reside on the same node as the lock manager.

As a result, *Wikstrom* cannot possible teach or suggest the feature of “**wherein each of the plurality of local lock manager processes may grant a lock on the same resource.**”

The Office Action argues that *Wikstrom* teaches this feature at Col. 6, lines 38-44; lines 54-64, and FIG. 5. However, this portion of *Wikstrom* merely describes sending a lock request message to each lock manager (LM) residing at each node. When each lock manager receives a lock request message, each lock manager locks a different resource, namely a particular version of a data object that resides at the same node as the lock manager. Nothing in the cited portion of *Wikstrom* teaches or suggests that two or more lock managers may grant locks on the same resource. If the Office disagrees, the Office is respectfully requested to specifically identify, within the teachings of *Wikstrom*, (a) the resource (and the location of the resource) that is the subject of the lock granted by multiple local lock managers, and (b) the two or more local lock managers that may each grant a lock on the same resource.

Consequently, *Wikstrom* also cannot teach a plurality of local lock manager processes that may each grant a lock on the same resource.

As a result of the fundamental differences between Claim 8 and the approach of *Troxel*, *Soltis*, and *Wikstrom*, numerous features of Claim 8 are not disclosed, taught, or suggested by *Troxel*, *Soltis*, or *Wikstrom*, either individually or in combination.

For example, Claim 8 recites the feature of “*providing a distributed lock manager process comprising a plurality of local lock manager processes executing on a corresponding plurality of hosts*” and “*wherein each of the plurality of local lock manager processes may grant a lock on the same resource.*” The single, centralized server of *Troxel* that is responsible for

locking resources upon receipt of a client request (Col. 1, lines 15-36) does not suggest a plurality of local lock manger processes, executing on a corresponding plurality of hosts, that each may grant a lock on the same resource. Indeed, the Office Action acknowledges, “*Troxel* does not explicitly teach providing a distributed lock manager process comprising a plurality of local lock manager processes executing on a corresponding plurality of hosts,” and instead, relies upon *Soltis* to show these features.

The portion of *Soltis* cited to show these features (FIG. 4; Clients 105A-N) merely shows a plurality of clients. Clients 105A-N of FIG. 4 request locks on data blocks from the data storage devices in the Network Storage Pool (NSP) 400. While each of the data storage devices, in the NSP 400, may grant a lock on a data block, a particular data storage device only grants locks to data blocks of that particular data storage device (see FIG. 7; Col. 13, line 65 – Col. Col. line 32). Thus, *Soltis* expressly teaches away from the feature of “wherein each of the plurality of local lock manager processes may grant a lock on the same resource.”

Moreover, *Wikstrom* cannot teach or suggest this feature as well, because in *Wikstrom*, each lock manager receives a lock request message, each lock manager locks a different resource, namely a particular version of a data object that resides at the same node as the lock manager. Thus, the above-quoted elements are not disclosed, taught, or suggested by *Troxel*, *Soltis*, or *Wikstrom*, either individually or in combination.

Claim 8 also recites the feature of “*associating a user identification for each user with one host of the plurality of hosts.*” The portion of *Soltis* cited to show this feature (FIG. 4; Clients 105A-N; FIGS. 9 and 10, Abstract, Col. 2, lines 47 to Col. 3, line 20) merely shows a plurality of clients. No explanation is provided by the Office Action as to why clients 105 A-N of FIG. 4 show associating a user identification for each user with one host of the plurality of

hosts. Importantly, the entire disclosure of *Soltis* lacks any teaching of associating a user identification for each user with any host. Further, each host, of the plurality of hosts, to which the user identifications are associated, is executing a local lock manager process. As explained above, *Soltis* lacks any teaching of a local lock manger process as claimed; consequently, it follows that *Soltis* lacks any teaching of associating a user identification for each user with one host of the plurality of hosts that are each executing a local lock manager process as claimed. Thus, this element cannot be disclosed, taught, or suggested by *Soltis*. The Office Action acknowledges that *Troxel* fails to teach this element, and does not cite *Wikstrom* to show this element.

Claim 8 also recites the feature of “*responding to a request for the resource associated with a first user having a first user identification associated with a first host of the plurality of hosts by requesting a lock from a first local lock manager process executing on the first host.*” *Soltis* is cited to show this element. However, as explained above, *Soltis* lacks any teaching or suggestion of (a) a local lock manager as claimed, (b) a user identification as claimed. Thus, *Soltis* cannot possibly teach or suggest this element, which involves a local lock manager and a user identification.

At best, the portion of *Soltis* cited to show this feature discusses that clients may acquire locks for shared use by multiple clients. However, *Soltis* teaches that a particular data storage device may only grant locks for data blocks of the particular data storage device (see FIG. 7; Col. 13, line 65 – Col. Col. line 32). Thus, when a data storage device receives a request for a resource that involves a lock, rather than requesting a lock from a local lock manager process executing on a host associated with the user identification that is associated with the requesting user, the data storage device processes the request without communicating with a local lock

manager process. As a result, this element is not disclosed, taught, or suggested by *Soltis*. The Office Action acknowledges that *Troxel* fails to teach this element, and does not cite *Wikstrom* to show this element.

As at least one element of Claim 8 is not disclosed, taught, or suggested by *Troxel*, *Soltis*, or *Wikstrom*, either individually or in combination, it is respectfully submitted that Claim 8 is patentable over the cited art and is in condition for allowance.

Moreover, *Troxel*, *Soltis*, and *Wikstrom* also have not been properly combined, and as a result, the rejection under 35 U.S.C. § 103(a) may not be properly maintained. The Office Action states that it would have been obvious to “combine the teachings of *Troxel* with the teaching of *Soltis* and *Wikstrom* in order to provide decentralized control of shared data, therefore maintaining data consistency¹,” but nothing in either *Troxel*, *Soltis*, or *Wikstrom* teaches or suggests combining their respective teachings².

As stated in the Federal Circuit decision *In re Dembiczak*, 50 USPQ.2d 1617 (Fed. Cir. 1999), (citing *Gore v. Garlock*, 220 USPQ 303, 313 (Fed. Cir. 1983)), “it is very easy to fall victim to the insidious effect of the hindsight syndrome where that which only the inventor taught is used against its teacher.” *Id.* The Federal Circuit stated in *Dembiczak* “that the best defense against subtle but powerful attraction of a hindsight-based obviousness analysis is rigorous application of the requirement for a showing of the teaching or suggestion to combine prior art references.” *Id.* Thus, the Federal Circuit explains that a proper obviousness analysis

¹ This is exactly the same motivation supplied previously, word for word, except now the motivation includes a third reference. This type of boilerplate motivation to combine is a textbook example of an impermissible hindsight-based rejection.

² The Office Action also completely fails to show how the approaches of *Troxel*, *Soltis*, and *Wikstrom* may possibly be used in conjunction with one another. For example, a single entity performs locking functions in *Troxel* and multiple entities perform locking functions in *Soltis* and *Wikstrom*. As another example of how it is unclear how the references may be combined, the resources being locked are different, e.g., in *Troxel*, files are locked, in *Soltis*, blocks are locked, and in *Wikstrom*, replicated data objects are locked.

requires “*particular factual findings* regarding the locus of the suggestion, teaching, or motivation to combine prior art references.” *Id.* (emphasis added).

In particular, the Federal Circuit states:

“We have noted that evidence of a suggestion, teaching, or motivation to combine may flow from the prior art references themselves, the knowledge of one of ordinary skill in the art, or, in some cases, from the nature of the problem to be solved...although ‘the suggestion more often comes from the teachings of the pertinent references’...The range of sources available, however, does *not diminish the requirement for actual evidence*. That is, the *showing must be clear and particular*...Broad conclusory statements regarding the teaching of multiple references, standing alone, are not ‘evidence.’” *Id.* (emphasis added; internal citations omitted).

Neither *Troxel*, nor *Soltis*, nor *Wikstrom* show any suggestion, teaching, or motivation to combine their teachings, nor does the Office Action provide a “clear and particular” showing of the suggestion, teaching, or motivation to combine their teachings. Nor portion of any reference is cited to support a motivation to combine. In fact, the only motivation provided in the Office Action is the hindsight observation that by combining features of those references, one may achieve the benefits achieved from the invention as described and claimed in the application. Such a hindsight observation is not consistent with the Federal Circuit’s requirement for “particular factual findings.”

CLAIMS 32, 37, AND 42

Claim 32 is a computer-readable medium claim that features limitations similar to those recited in method Claims 8. Consequently, it is respectfully submitted that Claim 32 is patentable over the cited art and are in condition for allowance for at least the reasons given above with respect to Claim 8.

Claim 37 is an apparatus claim in accordance with 35 U.S.C. § 112, sixth paragraph, which features limitations similar to those recited in method Claim 8. Consequently, it is

respectfully submitted that Claim 37 is patentable over the cited art and is in condition for allowance for at least the reasons given above with respect to Claim 8.

Claim 42 is an apparatus claim that features limitations similar to those recited in method Claim 8. Consequently, it is respectfully submitted that Claim 42 is patentable over the cited art and is in condition for allowance for at least the reasons given above with respect to Claim 8.

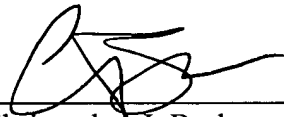
CONCLUSION

It is respectfully submitted that all of the pending claims are in condition for allowance and the issuance of a notice of allowance is respectfully requested. If there are any additional charges, please charge them to Deposit Account No. 50-1302.

The Examiner is invited to contact the undersigned by telephone if the Examiner believes that such contact would be helpful in furthering the prosecution of this application.

Respectfully submitted,

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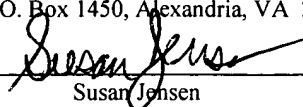
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on December 8, 2006

by


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